

A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney

Outline

# A Cybersecurity Investment Supply Chain Game Theory Model

#### Patrizia Daniele A. Maugeri A. Nagurney

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Session: Game Theory



## Cyber Attacks

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#### Verizon's 2016 Data Breach Investigations Report

2,260 confirmed data breaches at organizations in 82 countries

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## Recent Cyber Attacks

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Late summer 2014: 76 million customers



## Recent Cyber Attacks

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Late 2013: 40 million payment cards stolen and upwards of 70 million other personal records compromised

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## Recent Cyber Attacks

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Late Fall 2014: catastrophic and a public relations nightmare



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#### Center for Strategic and International Sudies (2014)

It has been estimated that the world economy sustained \$445 billion in losses from cyberattacks in 2014.

The estimated annual cost to the global economy from cybercrime is more than \$400 billion with a conservative estimate being \$375 billion in losses, a number that exceeds the national income of most countries.

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#### Pricewaterhouse Coopers (2014)

The number of cybersecurity incidents that were detected by respondents to their survey increased by 48% to 42.8 million in 2014



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The number of cybersecurity incidents that were detected by respondents to their survey increased by 48% to 42.8 million in 2014

No industrial sector is immune to cyber attacks with sectors such as financial services, insurance, pharmaceuticals, healthcare, high technology, energy, automotive and governments being especially attractive targets.

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#### Kaspersky Lab (2015)

A multinational gang of cybercriminals, known as *Carbanak*, infiltrated more than 100 banks across 30 countries and extracted as much as one billion dollars over a period of roughly two years

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#### Forbes (2015)

Cyberattacks can result not only in direct financial losses and/or the loss of data, but also in an organization's highly valued asset - its reputation

World-wide spending on cybersecurity was approximately \$75 billion in 2015, with the expectation that, by 2020, companies around the globe will be spending around \$170 billion annually



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Numerous companies and organizations have now realized that investing in cybersecurity is an imperative



# Some Bibliography

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- A. Nagurney, *Service Science* (2015): developed a multiproduct network economic model of cybercrime with a focus on financial services, since that industrial sector is a major target of cyberattacks
- A. Nagurney, L.S. Nagurney, *Netnomics* (2015): constructed a supply chain game theory model in which sellers maximize their expected profits while determining both their product transactions with consumers as well as their cybersecurity investments
- A. Nagurney, L.S. Nagurney, S. Shukla, in Computation, Cryptography, and Network Security (2015): extended the model to quantify and compute network vulnerability



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Patrizia Daniele, A. Maugeri, A. Nagurney

- A. Nagurney, P. D., S. Shukla, *Ann. Oper. Res.* (2017): introduced a novel game theory model in which the budget constraints for cybersecurity investments of retailers, which are nonlinear, are explicitly included, and conducted a spectrum of sensitivity analysis exercises
- P.D., A. Maugeri, A. Nagurney, in **Operations Research**, **Engineering**, and **Cyber Security** (2017): provided an alternative formulation of the variational inequality and a deeper qualitative and economic analysis with a focus on the Lagrange multipliers associated with the constraints

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Demand Markets

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Figure: The Bipartite Structure of the Supply Chain Network Game Theory Model



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#### Aim:

The retailers seek to maximize their individual expected utilities, consisting of expected profits, and compete in a noncooperative game in terms of strategies consisting of their respective product transactions and security levels

#### Conservation Law:

$$d_j = \sum_{i=1}^m Q_{ij}, \quad j = 1, \dots, n$$



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#### Upper Bounds for Production Transactions

$$0 \leq Q_{ij} \leq ar{Q}_{ij}, \hspace{1em} i=1,\ldots,m; j=1,\ldots,n$$

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A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney Upper Bounds for Production Transactions

$$0 \leq Q_{ij} \leq ar{Q}_{ij}, \quad i=1,\ldots,m; j=1,\ldots,n$$

Upper Bounds for Cybersecurity Levels

$$0 \leq s_i \leq u_{s_i}, \quad i = 1, \ldots, m, \quad \text{where } u_{s_i} < 1$$

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A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney Upper Bounds for Production Transactions

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Upper Bounds for Cybersecurity Levels

$$0 \leq s_i \leq u_{s_i}, \quad i = 1, \dots, m, \quad \text{where } u_{s_i} < 1$$

Demand Price of the Product at Demand Market j

$$\rho_j(d, \bar{s}) \equiv \hat{\rho}_j(Q, \bar{s}); \quad j = 1, \dots, n$$

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A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney Upper Bounds for Production Transactions

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Upper Bounds for Cybersecurity Levels

$$0 \leq s_i \leq u_{s_i}, \quad i = 1, \ldots, m, \quad \text{where } u_{s_i} < 1$$

Demand Price of the Product at Demand Market j

$$\rho_j(d,\bar{s}) \equiv \hat{\rho}_j(Q,\bar{s}); \quad j=1,\ldots,n$$

Investment Cost Function Associated with Achieving a Security Level  $s_i$ 

$$h_i(s_i) = lpha_i \left( rac{1}{\sqrt{1-s_i}} - 1 
ight)$$
 with  $lpha_i > 0$ 



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#### Budget Constraint

$$\alpha_i\left(\frac{1}{\sqrt{(1-s_i)}}-1\right)\leq B_i;\quad i=1,\ldots,m,$$

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#### Budget Constraint

$$lpha_i\left(rac{1}{\sqrt{(1-s_i)}}-1
ight)\leq B_i;\quad i=1,\ldots,m,$$

#### Profit of Retailer i

$$f_i(Q,s) = \sum_{j=1}^n \hat{
ho}_j(Q,s) Q_{ij} - c_i \sum_{j=1}^n Q_{ij} - \sum_{j=1}^n c_{ij}(Q_{ij})$$

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## Utility Optimization

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Each retailer seeks to maximize his expected utility

$$\max E(U_i) = (1 - p_i)f_i(Q, s) + p_i(f_i(Q, s) - D_i) - h_i(s_i)$$
$$= f_i(Q, s) - p_iD_i - h_i(s_i)$$

where:

- D<sub>i</sub> : damage incurred by retailer i
- $p_i = (1 s_i)(1 \bar{s}), \quad i = 1, ..., m$ : probability of a successful cyberattack on retailer *i*



# Nash Equilibrium

A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney Definition (A Supply Chain Nash Equilibrium in Product Transactions and Security Levels)

A product transaction and security level pattern  $(Q^*, s^*) \in \mathbb{K}$  is said to constitute a supply chain Nash equilibrium if for each retailer i; i = 1, ..., m,

$$\mathsf{E}(U_i(Q_i^*,s_i^*,\hat{Q}_i^*,\hat{s}_i^*)) \geq \mathsf{E}(U_i(Q_i,s_i,\hat{Q}_i^*,\hat{s}_i^*)), \quad orall (Q_i,s_i) \in \mathbb{K}^d$$

where

$$\hat{Q_i^*} \equiv (Q_1^*, \dots, Q_{i-1}^*, Q_{i+1}^*, \dots, Q_m^*);$$
 and  
 $\hat{s_i^*} \equiv (s_1^*, \dots, s_{i-1}^*, s_{i+1}^*, \dots, s_m^*)$ 



# Nash Equilibrium

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where

$$\hat{Q_i^*} \equiv (Q_1^*, \dots, Q_{i-1}^*, Q_{i+1}^*, \dots, Q_m^*);$$
 and  
 $\hat{s_i^*} \equiv (s_1^*, \dots, s_{i-1}^*, s_{i+1}^*, \dots, s_m^*)$ 

A supply chain Nash equilibrium is established if no retailer can unilaterally improve upon his expected utility (expected profit) by choosing an alternative vector of product transactions and security level.



# Variational Inequality Formulation

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#### Theorem

Assume that  $E(U_i(Q, s))$ , i = 1, ..., m is concave and continuously differentiable. Then  $(Q^*, s^*) \in \mathbb{K}$  is a supply chain Nash equilibrium  $\iff$  if it satisfies variational inequality

$$-\sum_{i=1}^m\sum_{j=1}^nrac{\partial \mathcal{E}(U_i(Q^*,s^*))}{\partial Q_{ij}} imes ig(Q_{ij}-Q_{ij}^*ig)$$

$$-\sum_{i=1}^m rac{\partial \mathcal{E}(U_i(Q^*,s^*))}{\partial s_i} imes (s_i-s_i^*) \geq 0, \quad orall (Q,s) \in \mathbb{K}$$



# Variational Inequality Formulation

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Feasible Set

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$$\mathbb{K} = \left\{ (Q, s) \in \mathbb{R}^{mn+n} : -Q_{ij} \leq 0, \ Q_{ij} - \overline{Q}_{ij} \leq 0, \ -s_i \leq 0, \\ s_i - u_{s_i} \leq 0, \ h_i(s_i) - B_i \leq 0, \ i = 1, \dots, m, \ j = 1, \dots, n \right\}$$



# Variational Inequality Formulation

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$$\begin{split} \mathbb{K} &= \left\{ (Q,s) \in \mathbb{R}^{mn+n} : -Q_{ij} \leq 0, \ Q_{ij} - \overline{Q}_{ij} \leq 0, \ -s_i \leq 0, \\ s_i - u_{s_i} \leq 0, \quad h_i(s_i) - B_i \leq 0, \ i = 1, \dots, m, \ j = 1, \dots, n \right\} \end{split}$$

#### Minimization Problem

Feasible Set

$$V(Q,s) \geq 0$$
 in  $\mathbb K$  and  $\min_{\mathbb K} V(Q,s) = V(Q^*,s^*) = 0$ , where

$$egin{aligned} \mathcal{V}(\mathcal{Q},s) &= & -\sum_{i=1}^m \sum_{j=1}^n rac{\partial \mathcal{E}(\mathcal{U}_i(\mathcal{Q}^*,s^*))}{\partial \mathcal{Q}_{ij}} \left(\mathcal{Q}_{ij}-\mathcal{Q}_{ij}^*
ight) \ & -\sum_{i=1}^m rac{\partial \mathcal{E}(\mathcal{U}_i(\mathcal{Q}^*,s^*))}{\partial s_i} \left(s_i-s_i^*
ight) \end{aligned}$$



## The Lagrange Theory

Lagrange Function

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# $\mathcal{L}(Q, s, \lambda^{1}, \lambda^{2}, \mu^{1}, \mu^{2}, \lambda) = -\sum_{i=1}^{m} \sum_{j=1}^{n} \frac{\partial E(U_{i}(Q^{*}, s^{*}))}{\partial Q_{ij}} (Q_{ij} - Q_{ij}^{*})$ $-\sum_{i=1}^{m} \frac{\partial E(U_{i}(Q^{*}, s^{*}))}{\partial s_{i}} (s_{i} - s_{i}^{*}) + \sum_{i=1}^{m} \sum_{j=1}^{n} \lambda_{ij}^{1} (-Q_{ij})$ $+\sum_{i=1}^{m} \sum_{j=1}^{n} \lambda_{ij}^{2} (Q_{ij} - \overline{Q}_{ij}) + \sum_{i=1}^{m} \mu_{i}^{1} (-s_{i})$

$$+\sum_{i=1}\mu_{i}^{2}(s_{i}-u_{s_{i}})+\sum_{i=1}\lambda_{i}(h_{i}(s_{i})-B_{i}),$$

where  $(Q, s) \in \mathbb{R}^{mn+n}, \lambda^1, \lambda^2 \in \mathbb{R}^{mn}_+, \mu^1, \mu^2 \in \mathbb{R}^m_+$ 

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# The Lagrange Theory

Theorem (Saddle Point)

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Patrizia Daniele, A. Maugeri, A. Nagurney There exist  $\overline{\lambda}^1$ ,  $\overline{\lambda}^2 \in \mathbb{R}^{mn}_+$ ,  $\overline{\mu}^1$ ,  $\overline{\mu}^2$ ,  $\overline{\lambda} \in \mathbb{R}^m_+$  such that the vector  $(Q^*, s^*, \overline{\lambda}^1, \overline{\lambda}^2, \overline{\mu}^1, \overline{\mu}^2, \overline{\lambda})$  is a saddle point of the Lagrange function; namely,

$$egin{aligned} \mathcal{L}(m{Q}^*,m{s}^*,\lambda^1,\lambda^2,\mu^1,\mu^2,\lambda) &\leq & \mathcal{L}(m{Q}^*,m{s}^*,\overline{\lambda}^1,\overline{\lambda}^2,\overline{\mu}^1,\overline{\mu}^2,\overline{\lambda}) \ &\leq & \mathcal{L}(m{Q},m{s},\overline{\lambda}^1,\overline{\lambda}^2,\overline{\mu}^1,\overline{\mu}^2,\overline{\lambda}) \end{aligned}$$

 $orall (\mathcal{Q}, s) \in \mathbb{K}, \, orall \lambda^1, \lambda^2 \in \mathbb{R}^{mn}_+, \, orall \mu^1, \mu^2, \lambda \in \mathbb{R}^m_+$  and

$$\overline{\lambda}_{ij}^1(-Q_{ij}^*)=0, \quad \overline{\lambda}_{ij}^2(Q_{ij}^*-\overline{Q}_{ij})=0, \quad orall i, \ orall j$$

 $\overline{\mu}_i^1(-s_i^*)=0, \quad \overline{\mu}_i^2(s_i^*-u_{s_i})=0, \quad \overline{\lambda}_i(h_i(s_i^*)-B_i)=0, \quad \forall i$ 

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## The Lagrange Theory

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Patrizia Daniele, A. Maugeri, A. Nagurney It follows that  $(Q^*, s^*) \in \mathbb{R}^{mn+n}_+$  is a minimal point of  $\mathcal{L}(Q, s, \overline{\lambda}^1, \overline{\lambda}^2, \overline{\mu}^1, \overline{\mu}^2, \overline{\lambda})$  in the whole space  $\mathbb{R}^{mn+n}$  and, hence, for all  $i = 1, \dots, m$ , and  $j = 1, \dots, n$ , we get:

$$\frac{\partial \mathcal{L}(Q^*, s^*, \overline{\lambda}^1, \overline{\lambda}^2, \overline{\mu}^1, \overline{\mu}^2, \overline{\lambda})}{\partial Q_{ij}} = -\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}} - \overline{\lambda}_{ij}^1 + \overline{\lambda}_{ij}^2 = 0$$
$$\frac{\partial \mathcal{L}(Q^*, s^*, \overline{\lambda}^1, \overline{\lambda}^2, \overline{\mu}^1, \overline{\mu}^2, \overline{\lambda})}{\partial s_i} = -\frac{\partial E(U_i(Q^*, s^*))}{\partial s_i}$$
$$-\overline{\mu}_i^1 + \overline{\mu}_i^2 + \overline{\lambda}_i \frac{\partial h_i(s^*_i)}{\partial s_i} = 0$$

which represent an equivalent formulation of the variational inequality



## Expected Utilities

A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney We define:

• 
$$\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}}$$
: the marginal expected transaction utility,  
 $i = 1, ..., m, j = 1, ..., n,$ 

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## **Expected Utilities**

A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney We define: •  $\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}}$ : the marginal expected transaction utility,  $i = 1, \dots, m, j = 1, \dots, n,$ •  $\frac{\partial E(U_i(Q^*, s^*))}{\partial s_i}$ : the marginal expected cybersecurity investment utility,  $i = 1, \dots, m$ 



## **Expected Utilities**

A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney We define: •  $\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}}$ : the marginal expected transaction utility, i = 1, ..., m, j = 1, ..., n,•  $\frac{\partial E(U_i(Q^*, s^*))}{\partial s_i}$ : the marginal expected cybersecurity investment utility, i = 1, ..., m $\overline{r^1}$   $\overline{r^2}$ 

 $\overline{\lambda}_{ij}^1, \overline{\lambda}_{ij}^2$  give a precise evaluation of the behavior of the market with respect to the supply chain product transactions as well as  $\overline{\mu}_i^1, \overline{\mu}_i^2$  describe the effects of the marginal expected cybersecurity investment utilities.



# Analysis of Marginal Expected Transaction Utilities

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$$-\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}} - \overline{\lambda}_{ij}^1 + \overline{\lambda}_{ij}^2 = 0, \quad i = 1, \dots, m, \ j = 1, \dots, n.$$
  
So, if  $0 < Q_{ij}^* < \overline{Q}_{ij}$ , then we get  $\forall i = 1, \dots, m, \ j = 1, \dots, n$ :  
$$-\frac{\partial E(U_i(Q^*, s^*))}{\partial Q_{ij}} = c_i + \frac{\partial c_{ij}(Q_{ij}^*)}{\partial Q_{ij}} - \hat{\rho}_j(Q^*, s^*) - \sum_{k=1}^m \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q_{ik}^* = 0,$$

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# Analysis of Marginal Expected Transaction Utilities

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We get  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{i:}} - \overline{\lambda}_{ij}^1 + \overline{\lambda}_{ij}^2 = 0, \quad i = 1, \dots, m, \ j = 1, \dots, n.$ So, if  $0 < Q_{ii}^* < \overline{Q}_{ij}$ , then we get  $\forall i = 1, \dots, m, j = 1, \dots, n$ :  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{ii}} = c_i + \frac{\partial c_{ij}(Q^*_{ij})}{\partial Q_{ii}} - \hat{\rho}_j(Q^*,s^*) - \sum_{i=1}^{m} \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q^*_{ik} = 0,$ whereas if  $\overline{\lambda}_{ij}^1 > 0$ , and, hence,  $Q_{ij}^* = 0$ , and  $\overline{\lambda}_{ij}^2 = 0$ , we get  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{ii}} = c_i + \frac{\partial c_{ij}(Q_{ij}^*)}{\partial Q_{ii}} - \hat{\rho}_j(Q^*,s^*) - \sum_{k=1}^{ii} \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q_{ik}^* = \overline{\lambda}_{ij}^1,$ 

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# Analysis of Marginal Expected Transaction Utilities

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We get  $-\frac{\partial E(U_i(Q^*, s^*))}{\partial \Omega_{::}} - \overline{\lambda}_{ij}^1 + \overline{\lambda}_{ij}^2 = 0, \quad i = 1, \dots, m, \ j = 1, \dots, n.$ So, if  $0 < Q_{ii}^* < \overline{Q}_{ij}$ , then we get  $\forall i = 1, \dots, m, j = 1, \dots, n$ :  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{ii}} = c_i + \frac{\partial c_{ij}(Q_{ij}^*)}{\partial Q_{ii}} - \hat{\rho}_j(Q^*,s^*) - \sum_{i=1}^{m} \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q_{ik}^* = 0,$ whereas if  $\overline{\lambda}_{ii}^1 > 0$ , and, hence,  $Q_{ii}^* = 0$ , and  $\overline{\lambda}_{ii}^2 = 0$ , we get  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{ij}} = c_i + \frac{\partial c_{ij}(Q^*_{ij})}{\partial Q_{ij}} - \hat{\rho}_j(Q^*,s^*) - \sum_{k=1}^{\prime\prime\prime} \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q^*_{ik} = \overline{\lambda}^1_{ij},$ and if  $\overline{\lambda}_{ii}^2 > 0$ , and, hence,  $Q_{ii}^* = \overline{Q}_{ij}$ , and  $\overline{\lambda}_{ii}^1 = 0$ , we have  $-\frac{\partial E(U_i(Q^*,s^*))}{\partial Q_{ii}} = c_i + \frac{\partial c_{ij}(Q^*_{ij})}{\partial Q_{ii}} - \hat{\rho}_j(Q^*,s^*) - \sum_{l=1}^{\prime\prime\prime} \frac{\partial \hat{\rho}_k}{\partial Q_{ij}} \times Q^*_{ik} = -\overline{\lambda}^2_{ij},$ < □ > < ₩≠i> < ≡ > < ≡ >



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We have 
$$\forall i = 1, \ldots, m$$
:

$$-\frac{\partial E(U_i(Q^*,s^*))}{\partial s_i} - \overline{\mu}_i^1 + \overline{\mu}_i^2 + \overline{\lambda}_i \frac{\partial h_i(s^*)}{\partial s_i} = 0,$$

If  $0 < s^*_i < u_{s_i},$  then  $\overline{\mu}^1_i = \overline{\mu}^2_i = 0$  and we have

$$\begin{aligned} & \frac{\partial h_i(s_i^*)}{\partial s_i} + \overline{\lambda}_i \frac{\partial h_i(s_i^*)}{\partial s_i} \\ & = \left(1 - \sum_{k=1}^m \frac{s_k^*}{m} + \frac{1 - s_i^*}{m}\right) D_i + \sum_{k=1}^m \frac{\partial \hat{\rho}_k(Q^*, s^*)}{\partial s_i} \times Q_{ik}^*. \end{aligned}$$



A Cybersecurity Investment Supply Chain Game Theory Model

Patrizia Daniele, A. Maugeri, A. Nagurney We have  $\forall i = 1, \dots, m$ :

$$-\frac{\partial E(U_i(Q^*,s^*))}{\partial s_i} - \overline{\mu}_i^1 + \overline{\mu}_i^2 + \overline{\lambda}_i \frac{\partial h_i(s^*)}{\partial s_i} = 0,$$

If  $0 < s^*_i < u_{s_i},$  then  $\overline{\mu}^1_i = \overline{\mu}^2_i = 0$  and we have

$$\begin{aligned} \frac{\partial h_i(\boldsymbol{s}_i^*)}{\partial \boldsymbol{s}_i} + \overline{\lambda}_i \frac{\partial h_i(\boldsymbol{s}_i^*)}{\partial \boldsymbol{s}_i} \\ = \left(1 - \sum_{k=1}^m \frac{\boldsymbol{s}_k^*}{m} + \frac{1 - \boldsymbol{s}_i^*}{m}\right) D_i + \sum_{k=1}^m \frac{\partial \hat{\rho}_k(\boldsymbol{Q}^*, \boldsymbol{s}^*)}{\partial \boldsymbol{s}_i} \times \boldsymbol{Q}_{ik}^*. \end{aligned}$$

Since  $0 < s_i^* < u_{s_i}$ ,  $h(s_i^*)$  cannot be the upper bound  $B_i$ ; hence,  $\overline{\lambda}_i$  is zero and hence

$$\frac{\partial h_i(s_i^*)}{\partial s_i} = \left(1 - \sum_{k=1}^m \frac{s_k^*}{m} + \frac{1 - s_i^*}{m}\right) D_i + \sum_{k=1}^m \frac{\partial \hat{\rho}_k(Q^*, s^*)}{\partial s_i} \times Q_{ik}^*$$



A Cybersecurity Investment Supply Chain Game Theory Model

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If 
$$\overline{\mu}_i^1 > 0$$
 and, hence,  $s_i^* = 0$ , and  $\overline{\mu}_i^2 = 0$ , we get:  

$$-\frac{\partial E(U_i(Q^*, s^*))}{\partial s_i}$$

$$= \frac{\partial h_i(0)}{\partial s_i} - \left(1 - \sum_{k=1}^m \frac{s_k^*}{m} + \frac{1 - s_i^*}{m}\right) D_i - \sum_{k=1}^m \frac{\partial \hat{\rho}_k(Q^*, s^*)}{\partial s_i} Q_{ik}^* = \overline{\mu}_i^1.$$



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In contrast, if  $\overline{\mu}_i^2 > 0$  and, hence,  $s_i^* = u_{s_i}$ , retailer j has a marginal gain given by  $\overline{\mu}_i^2$ , because

$$-\frac{\partial E(U_i(Q^*, u_{s_i}))}{\partial s_i} = -\left(1 - \sum_{\substack{k=1\\k \neq i}}^m \frac{u_{s_k}}{m} + \frac{1 - u_{s_i}}{m}\right) D_i$$
$$-\sum_{k=1}^m \frac{\partial \hat{\rho}_k(Q^*, s^*)}{\partial s_i} \times Q_{ik}^* + \frac{\partial h_i(u_{s_i})}{\partial s_i} + \overline{\lambda}_i \frac{\partial h_i(u_{s_i})}{\partial s_i} = -\overline{\mu}_i^2.$$





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Demand Markets

Figure: Network Topology

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#### Cost Functions

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#### **Demand Price Functions**

Cost Functions

$$\rho_1(d,\overline{s}) = -d_1 + .1 \frac{s_1 + s_2}{2} + 100, \quad \rho_2(d,\overline{s}) = -.5d_2 + .2 \frac{s_1 + s_2}{2} + 200$$

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#### **Demand Price Functions**

**Cost Functions** 

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#### Damage Parameters and Budgets

 $D_1 = 200$  and  $D_2 = 210$ ,  $B_1 = B_2 = 2.5$  in millions of \$



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#### Equilibrium Solution

$$Q_{11}^* = 24.148$$
,  $Q_{21}^* = 21.586$ ,  $Q_{12}^* = 99.16$ ,  $Q_{22}^* = 94.16$ ,  
 $\overline{\mu}_1^2 = 19.6055$ ,  $\overline{\mu}_2^2 = 20.3273$ ,  
where  $\overline{\mu}_1^2$  and  $\mu_2^2$  are the positive marginal expected gains.

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#### Equilibrium Solution

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here  $\overline{\mu}_1^2$  and  $\mu_2^2$  are the positive marginal expected gains.

If we double the value of the damage for the first retailer and assume now  $D_1 = 400$ , then the new value of the Lagrange multiplier is  $\overline{\mu}_1^2 = 46.6055$ .

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- Cyberattacks are negatively globally impacting numerous sectors of economies

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• Organizations are investing in cybersecurity



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- Cyberattacks are negatively globally impacting numerous sectors of economies
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- Retailers compete in both product transactions and cybersecurity levels seeking to maximize their expected utilities



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- Organizations are investing in cybersecurity
- Retailers compete in both product transactions and cybersecurity levels seeking to maximize their expected utilities
- The governing equilibrium concept is that of Nash equilibrium
- We perform an analysis of both the marginal expected transaction utilities and the marginal expected cybersecurity investment utilities of the retailers

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